State transition models of biomolecular dynamics A few research questions

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 - Research questions: metabolic P system models

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State transition dynamics: basic concepts

state transition (ST) dynamics: (S, q)

- S: (discrete) state space
- $q: S{
 ightarrow} \mathcal{P}(S)$: [total] transition **dynamics**
- $\mathbf{x} \rightarrow \mathbf{y} \stackrel{\text{def}}{=} \mathbf{y} \in \mathbf{q}(\mathbf{x})$: (binary) transition relation
- $q(X) \stackrel{\text{def}}{=} \bigcup_{s \in X} q(s)$: extension to **quasistates** $X \subseteq S$
- orbit : $(X_i | i \in \mathbb{N})$ t.c. $\forall i. X_{i+1} = q(X_i)$
- eventually periodic orbit of origin X : $\exists k \ge 0 \exists n > 0. q^{k+n}(X) = q^k(X)$ (periodic orbit if k = 0)
- orbit (eventually) included in another one: $(\exists i \ge 0.) \forall j (\ge i) X_j \subseteq X'_i$
- **basin** $B : \emptyset \neq B \subseteq S$ s.t. $q(B) \subseteq B$
- s-trajectory : ξ : $\mathbb{N} \rightarrow S$ s.t. $\xi(0) = s$ and $\forall i \ge 0$. $\xi(i) \rightarrow \xi(i+1)$
- s-flight ^{def} injective s-trajectory

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orbit: (X_i|i ∈ N) t.c. ∀i. X_{i+1} = q(X_i)
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Basic concepts Recurrence and attractors Simple example

Deterministic ST dynamics: recurrence, attractors

Classical dynamics concepts:

attracting set, recurrence, attractor

- take two distinct modal flavours in the general, nondeterministic case, see [(Manca *et al.*, 2005), (Scollo *et al.*, 2006)]
- no modal difference in the deterministic case (where the concepts of orbit and trajectory practically coincide):
- state x is **recurrent** if $x \in q^n(x)$ for some n > 0
- an **attracting set** *A* of a basin *B* is a nonempty subset of *B* s.t. $\forall x \in B \exists k \in \mathbb{N} : \forall n \geq k \ q^n(x) \subseteq A$
- the **attractor** of a basin *B* is a minimal attracting set of *B* (under inclusion)
- the attractor of a basin B exists iff B has no flights
- when the attractor of *B* exists, it is the set of recurrent states in *B*

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ST dynamics: a contrived example

a simple model of epidemic propagation

as well as of unstable catalythic reaction:

- $CG \rightarrow GG$
- $C \rightarrow C$
- $G \rightarrow X$
- $G \rightarrow K$
- $\mathbf{G} \rightarrow \mathbf{G}$
- instability of agent G: it either dies or recovers (becoming immune: K)
- states : (|*C*|+|*K*|, |*G*|)
- transitions can be determined by equipping rules with a measure of relative strength

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- $\mathbf{G} \rightarrow \mathbf{K}$
- $\bm{G} \to \bm{G}$

• instability of agent G: it either dies or recovers (becoming immune: K)

• states : (|C| + |K|, |G|)

 transitions can be determined by equipping rules with a measure of relative strength

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Basic concepts Recurrence and attractors Simple example

ST dynamics: a contrived example

a simple model of epidemic propagation

as well as of unstable catalythic reaction:

 $\textbf{CG} \rightarrow \textbf{GG}$

- $\mathbf{C} \rightarrow \mathbf{C}$
- ${f G}
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P systems: basic concepts

P systems: a computational model based on

- multiset rewriting rules
- workspace compartmentalization
- biological counterparts:
 - biochemical reactions
 - membranes
- evolution strategies in (variants of) P systems:
 - maximal parallelism, usually
 - or some other fixed strategy
- this is inadequate to model most biological phenomena: evolution depends on the state, so strategy may change

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Metabolic P systems

a possible way out: **metabolic** P systems [(Bianco *et al.*, 2006a), (Bianco *et al.*, 2006b), (Manca, 2006)]:

metabolic rules: multiset rewriting rules equipped with reaction maps:

real-valued functions of the system state, that tell the **competitive strength** of each metabolic rule

- each metabolic rule:
 - consumes finite amounts of those biochemical resources whose types occur in the pattern (the lhs)
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generalizing Avogadro's molar units:

- only finite amounts of biochemical resources are available: their allocation to rules obeys a mass partition principle, where each rule gets an amount of each needed resource that is an integer multiple of the rule's reaction unit
- the latter depends on the reaction map and on the number of occurrences of resource types in the pattern
- **PMA** [(Manca, 2006)] is the simple algorithm that, at each state, computes the resource allocation to metabolic rules according to their reaction maps and reaction units

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Research questions: state transition dynamics Research questions: metabolic P system models

Research questions: state transition dynamics

Moving from structureless to topological state spaces:

definability of weaker notions of recurrence

that is: replace exact occurrence of a state in its own trajectory with approximate occurrence

definability of weaker notions of attraction

that is: replace exact inclusion of orbits in the attracting set with approximate inclusion

 generalization of the characterization results obtained in [(Scollo *et al.*, 2006)], linking recurrence and attractors, under the aforementioned weakings

at least for the deterministic case

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two levels in metabolic P system modeling of biochemical dynamics:

- stoichiometry: the metabolic rules
- regulation: the reaction maps
- basic question (roughly): how to get those models...

. . . for a given biochemical behaviour? This is the general **simulation problem**

easier (?) question: assume the rules are given, how to get their reaction maps?

This may be called the **regulation problem**

even easier (??) question: assumes the rules are given and the form of reaction maps is nown (polynomial, say), how to get the values of their unknown parameters (the coefficients, say)?

This may be called the tuning problem

an instructive case study:
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State transition dynamics Metabolic P systems

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