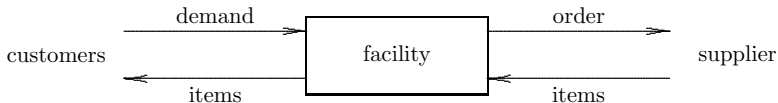


## Section 1.3: A Simple Inventory System

### Discrete-Event Simulation: A First Course

©2006 Pearson Ed., Inc. 0-13-142917-5

## Section 1.3: A Simple Inventory System



- Distributes items from current inventory to customers
- Customer demand is discrete
- Simple  $\iff$  one type of item

- *Transaction Reporting*
  - Inventory review after *each* transaction
  - Significant labor may be required
  - Less likely to experience shortage
- *Periodic Inventory Review*
  - Inventory review is periodic
  - Items are ordered, if necessary, only at review times
  - $(s, S)$  are the min,max inventory levels,  $0 \leq s < S$
- We assume periodic inventory review
- Search for  $(s, S)$  that minimize cost

## Inventory System Costs

- *Holding cost:* for items in inventory
- *Shortage cost:* for unmet demand
- *Setup cost:* fixed cost when order is placed
- *Item cost:* per-item order cost
- *Ordering cost:* sum of setup and item costs

## Additional Assumptions

- *Back ordering is possible*
- *No delivery lag*
- *Initial inventory level is  $S$*
- *Terminal inventory level is  $S$*

# Specification Model

- Time begins at  $t = 0$
- Review times are  $t = 0, 1, 2, \dots$
- $I_{i-1}$ : inventory level at *beginning* of  $i^{\text{th}}$  interval
- $o_{i-1}$ : amount ordered at time  $t = i - 1$ , ( $o_{i-1} \geq 0$ )
- $d_i$ : demand quantity *during*  $i^{\text{th}}$  interval, ( $d_i \geq 0$ )
- Inventory at end of interval can be negative

# Inventory Level Considerations

- Inventory level is reviewed at  $t = i - 1$
- If at least  $s$ , no order is placed  
If less than  $s$ , inventory is replenished to  $S$

$$o_{i-1} = \begin{cases} 0 & l_{i-1} \geq s \\ S - l_{i-1} & l_{i-1} < s \end{cases}$$

- Items are delivered immediately
- At end of  $i^{\text{th}}$  interval, inventory diminished by  $d_i$

$$l_i = l_{i-1} + o_{i-1} - d_i$$

# Time Evolution of Inventory Level

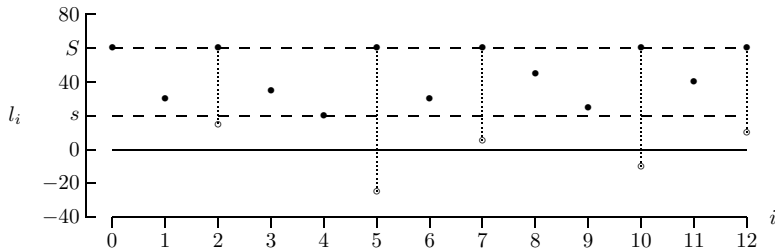
## Algorithm 1.3.1

```
 $l_0 = S; \quad /* \text{ the initial inventory level is } S */$   
 $i = 0;$   
while ( more demand to process ) {  
     $i++;$   
    if ( $l_{i-1} < s$ )  
         $o_{i-1} = S - l_{i-1};$   
    else  
         $o_{i-1} = 0;$   
     $d_i = \text{GetDemand}();$   
     $l_i = l_{i-1} + o_{i-1} - d_i;$   
}  
 $n = i;$   
 $o_n = S - l_n;$   
 $l_n = S; \quad /* \text{ the terminal inventory level is } S */$   
return  $l_1, l_2, \dots, l_n$  and  $o_1, o_2, \dots, o_n;$ 
```

# Example 1.3.1: SIS with Sample Demands

Let  $(s, S) = (20, 60)$  and consider  $n = 12$  time intervals

$i$	:	1	2	3	4	5	6	7	8	9	10	11	12
$d_i$	:	30	15	25	15	45	30	25	15	20	35	20	30





- What statistics to compute?
- *Average demand* and *average order*

$$\bar{d} = \frac{1}{n} \sum_{i=1}^n d_i$$

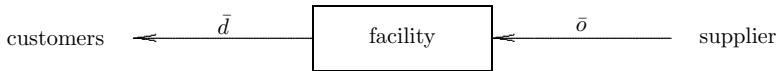
$$\bar{o} = \frac{1}{n} \sum_{i=1}^n o_i.$$

- For Example 1.3.1 data

$$\bar{d} = \bar{o} = 305/12 \simeq 25.42 \text{ items per time interval}$$

# Flow Balance

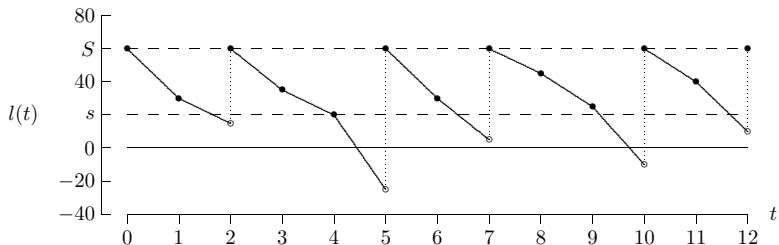
- Average demand and order *must* be equal
- Ending inventory level is  $S$
- Over the simulated period, all demand is satisfied
- Average “flow” of items in equals average “flow” of items out



- The inventory system is *flow balanced*

# Constant Demand Rate

- Holding and shortage costs are proportional to time-averaged inventory levels
- Must know inventory level for all  $t$
- *Assume* the demand rate is constant between review times

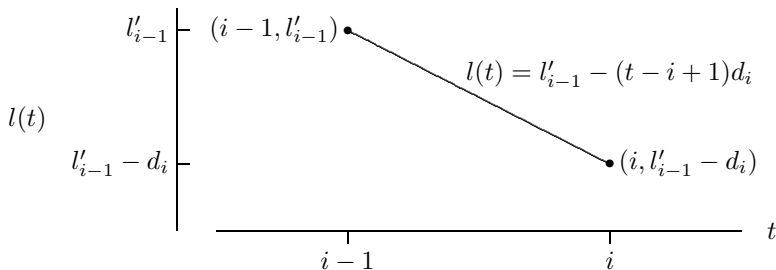


# Inventory Level as a Function of Time

- The inventory level at any time  $t$  in  $i^{\text{th}}$  interval is

$$l(t) = l'_{i-1} - (t - i + 1)d_i$$

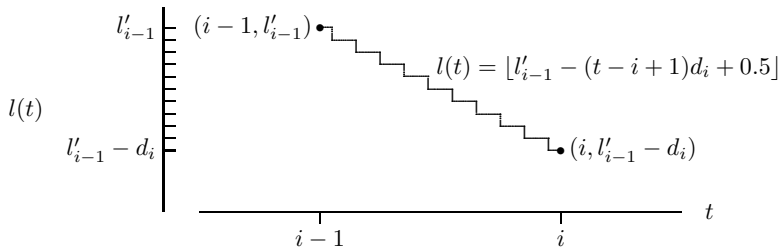
if demand rate is constant between review times



- $l'_{i-1} = l_{i-1} + o_{i-1}$  represents inventory level *after* review

# Inventory Decrease Is Not Linear

- Inventory level at any time  $t$  is an integer
- $I(t)$  should be rounded to an integer value
- $I(t)$  is a stair-step, rather than linear, function



# Time-Averaged Inventory Level

- $I(t)$  is the basis for computing the time-averaged inventory level
- Case 1: If  $I(t)$  remains non-negative over  $i^{\text{th}}$  interval

$$\bar{I}_i^+ = \int_{i-1}^i I(t) dt$$

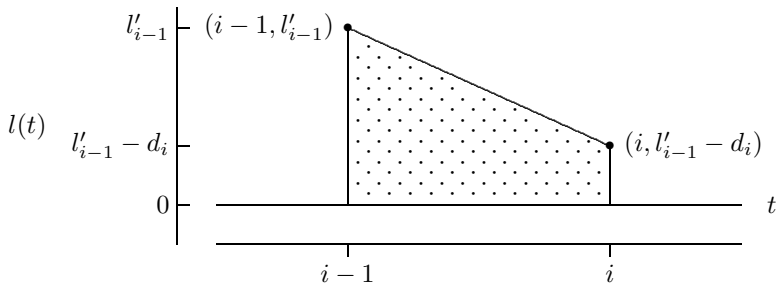
- Case 2: If  $I(t)$  becomes negative at some time  $\tau$

$$\bar{I}_i^+ = \int_{i-1}^{\tau} I(t) dt \qquad \bar{I}_i^- = - \int_{\tau}^i I(t) dt$$

- $\bar{I}_i^+$  is the time-averaged *holding level*  
 $\bar{I}_i^-$  is the time-averaged *shortage level*

# Case 1: No Back Ordering

- No shortage during  $i^{\text{th}}$  time interval iff.  $d_i \leq l'_{i-1}$

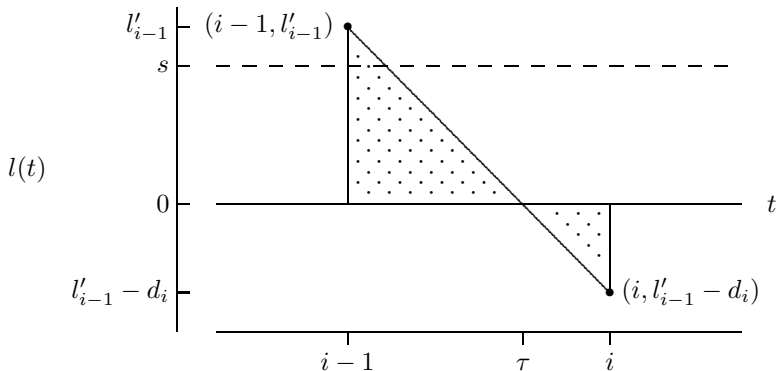


- Time-averaged holding level: area of a trapezoid

$$\bar{l}_i^+ = \int_{i-1}^i l(t) dt = \frac{l'_{i-1} + (l'_{i-1} - d_i)}{2} = l'_{i-1} - \frac{1}{2}d_i$$

## Case 2: Back Ordering

- Inventory becomes negative iff.  $d_i > l'_{i-1}$





## Case 2: Back Ordering (Cont.)

- $I(t)$  becomes negative at time  $t = \tau = i - 1 + (l'_{i-1}/d_i)$
- Time-averaged holding and shortage levels for  $i^{\text{th}}$  interval computed as the areas of triangles

$$\bar{l}_i^+ = \int_{i-1}^{\tau} I(t) dt = \dots = \frac{(l'_{i-1})^2}{2d_i}$$

$$\bar{l}_i^- = - \int_{\tau}^i I(t) dt = \dots = \frac{(d_i - l'_{i-1})^2}{2d_i}$$

# Time-Averaged Inventory Level

- *Time-averaged holding level and time-averaged shortage level*

$$\bar{l}^+ = \frac{1}{n} \sum_{i=1}^n \bar{l}_i^+ \qquad \bar{l}^- = \frac{1}{n} \sum_{i=1}^n \bar{l}_i^-$$

- Note that time-averaged shortage level is positive
- The *time-averaged inventory level* is

$$\bar{l} = \frac{1}{n} \int_0^n l(t) dt = \bar{l}^+ - \bar{l}^-$$

# Computational Model

- `sis1` is a trace-driven computational model of the SIS
- Computes the statistics

$$\bar{d}, \bar{o}, \bar{l}^+, \bar{l}^-$$

and the order frequency  $\bar{u}$

$$\bar{u} = \frac{\text{number of orders}}{n}$$

- Consistency check: compute  $\bar{o}$  and  $\bar{d}$  separately, then compare

## Example 1.3.4: Executing `sis1`

- Trace file `sis1.dat` contains data for  $n = 100$  time intervals
- With  $(s, S) = (20, 80)$

$$\bar{o} = \bar{d} = 29.29 \quad \bar{u} = 0.39 \quad \bar{T}^+ = 42.40 \quad \bar{T}^- = 0.25$$

- After Chapter 2, we will generate data randomly (no trace file)

A facility's cost of operation is determined by:

- $c_{\text{item}}$  : *unit cost of new item*
- $c_{\text{setup}}$  : *fixed cost for placing an order*
- $c_{\text{hold}}$  : *cost to hold one item for one time interval*
- $c_{\text{short}}$  : *cost of being short one item for one time interval*

- Automobile dealership that uses weekly periodic inventory review
- The facility is the showroom and surrounding areas
- The items are new cars
- The supplier is the car manufacturer
- “...customers are people convinced by clever advertising that their lives will be improved significantly if they purchase a new car from this dealer.” (S. Park)
- Simple (one type of car) inventory system

## Example 1.3.5: Case Study Materialized

- Limited to a maximum of  $S = 80$  cars
- Inventory reviewed every Monday
- If inventory falls below  $s = 20$ , order cars sufficient to restore to  $S$
- For now, ignore delivery lag
- Costs:
  - Item cost is  $C_{\text{item}} = \$8000$  per item
  - Setup cost is  $C_{\text{setup}} = \$1000$
  - Holding cost is  $C_{\text{hold}} = \$25$  per week
  - Shortage cost is  $C_{\text{hold}} = \$700$  per week

# Per-Interval Average Operating Costs

- The average operating costs *per time interval* are

- *item cost* :  $c_{\text{item}} \cdot \bar{d}$
- *setup cost* :  $c_{\text{setup}} \cdot \bar{u}$
- *holding cost* :  $c_{\text{hold}} \cdot \bar{I}^+$
- *shortage cost* :  $c_{\text{short}} \cdot \bar{I}^-$

- The average *total* operating cost *per time interval* is their sum
- For the stats and costs of the hypothetical dealership:

- *item cost* :  $\$8000 \cdot 29.29 = \$234,320$  *per week*
- *setup cost* :  $\$1000 \cdot 0.39 = \$390$  *per week*
- *holding cost* :  $\$25 \cdot 42.40 = \$1,060$  *per week*
- *shortage cost* :  $\$700 \cdot 0.25 = \$175$  *per week*



# Cost Minimization

- By varying  $s$  (and possibly  $S$ ), an optimal policy can be determined
- Optimal  $\iff$  minimum average cost
- Note that  $\bar{o} = \bar{d}$ , and  $\bar{d}$  depends only on the demands
- Hence, item cost is independent of  $(s, S)$
- Average *dependent* cost is  
avg setup cost + avg holding cost + avg shortage cost

- Let  $S$  be fixed, and let the demand sequence be fixed
- If  $s$  is systematically increased, we expect:
  - average setup cost and holding cost will increase as  $s$  increases
  - average shortage cost will decrease as  $s$  increases
  - average dependent cost will have 'U' shape, yielding an optimum
- From results (next slide), minimum cost is \$1550 at  $s = 22$

# Example 1.3.7: Simulation Results

