

Formal Analysis of Security Policies

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Security Workshop

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Development of an Access Control System

- 1) Authentication**
- 2) Security Policies**
- 3) Security Mechanisms**

Security Policy

**“A set of norms regulating the modalities
– obligation, permission, interdiction –
for a set of agents on some action ”**

Inconsistencies

- **Contradiction:**
“forbidden smoke” and “obligatory smoke”!
- **Dilemma:**
“forbidden smoke” and “forbidden no smoke”!



Example Policy



11 Norms
4 Roles

N1: `if play(a,User) and public(f)`
`then Perm(Read(a,f))`

N2: `if play(a,User) and public(f) and owner(f,a)`
`then Perm(Write(a,f))`

N3: `if play(a,User)`
`then Forb(Downgrade(a,f))`

N4: `if play(a,User) and password(a,p) and old(p)`
`then Obl(Change_Psswd(a))`

Example Policy

N5: `if play(a, Secret) and not(public(f))
then Perm(Read(a, f))`

N6: `if play(a, Secret) and not(public(f))
and owner(f, a)
then Perm(Write(a, f))`

N7: `if play(a, Sso)
then Perm(Downgrade(a, f))`

Example Policy

N8: if play(a,Bad)
 then Forb(Read(a,f))

N9: if play(a,Bad)
 then Forb(Write(a,f))

N10: if play(a,Bad)
 then Forb(Downgrade(a,f))

N11: if play(a,Bad)
 then Forb(Change_Psswd(a))

Inductive Approach

- **Trace:** list of admissible norms induced by policy
- **Model of Policy:** set of all possible trace of norms that the policy admits
 - Mechanized with the proof assistant:
PVS or Isabelle
- Properties of the model proved with the correspondent inductive principle

Inductive Definition of Policy

“Set of all possible trace of norms that the policy admits”

Base case

`[] ∈ Policy`

Inductive case

`trace ∈ Policy ⇒ nm # trace ∈ Policy`

Types

```
typedefcl Agent
```

```
typedefcl File
```

```
typedefcl Psswd
```

```
datatype Role = User | Sso | Secret | Bad
```

Functions

consts

play :: "[Agent, Role] ⇒ bool"

owner :: "[File, Agent] ⇒ bool"

password :: "[Agent, Psswd] ⇒ bool"

public :: "File ⇒ bool"

old :: "Psswd ⇒ bool"

Constraints on Roles

axioms

Secret_User [simp] : "play a Secret \rightarrow play a User"

Sso_Secret [simp] : "play a Sso \rightarrow play a Secret"

Bad_User [simp] : "play a Bad \rightarrow play a User"

Lemma Transitivity_Sso_User [simp] :

" \forall (a::Agent). play a Sso \rightarrow play a User"

Operations

```
datatype operation =
```

```
    Read          Agent File
```

```
| Write          Agent File
```

```
| Change_Psswd  Agent
```

```
| Downgrade     Agent File
```

```
| Not_op        operation    ("¬o")
```

```
axioms
```

```
Not_op_idemp [simp] : "¬o (¬o oper) = oper"
```

Norms

```
datatype norma =
```

```
    Obl          operation
  | Perm         operation
  | Forb         operation
  | Waived       operation
  | Not_norma   norma      ("¬n")
```

Axioms for Norms

axioms

Not_norma_idemp [simp]: " $\neg n (\neg n nm) = nm$ "

Perm_Obl [simp]: "Perm oper = $\neg n$ (Obl ($\neg o$ oper))"

Forb_Obl [simp]: "Forb oper = Obl ($\neg o$ oper)"

Mechanization with Isabelle

```
types trace = "norma list"

consts Policy :: "trace set"

inductive "Policy"

intros

Empty : "[ ] ∈ Policy"

Norma_1 : "[ | tr1 ∈ Policy; play a User;
           public f | ]
           ⇒ Perm (Read a f) # tr1 ∈ Policy"
```


Inconsistencies

Contradiction

$(\text{Obligatory}(\text{op}) \wedge \neg \text{Obligatory}(\text{op}))$

\vee

$(\text{Obligatory}(\neg \text{op}) \wedge \neg \text{Obligatory}(\neg \text{op}))$

Dilemma

$\text{Obligatory}(\text{op}) \wedge \text{Obligatory}(\neg \text{op})$

\vee

~~$(\neg \text{Obligatory}(\text{op}) \wedge \neg \text{Obligatory}(\neg \text{op}))$~~

Contradiction in Isabelle

```
consts Contradiction :: "norma  $\Rightarrow$  norma"
```

```
axioms Contradiction_1 [simp] :
```

```
"Contradiction (Obl oper) =  $\neg$ n (Obl oper)"
```

```
axioms Contradiction_2 [simp] :
```

```
"Contradiction ( $\neg$ n (Obl oper)) = Obl oper"
```

Dilemma in Isabelle

```
consts Dilemma :: "norma  $\Rightarrow$  norma"
```

```
axioms Dilemma_1 [simp] :
```

```
"Dilemma (Obl oper) = Obl ( $\neg$ o oper)"
```

Absence of Contradictions and of Dilemmas

theorem No_Contradiction :

“ [| nm ∈ set tr; tr ∈ Policy |] ⇒
Contradiction nm ∉ set tr ”

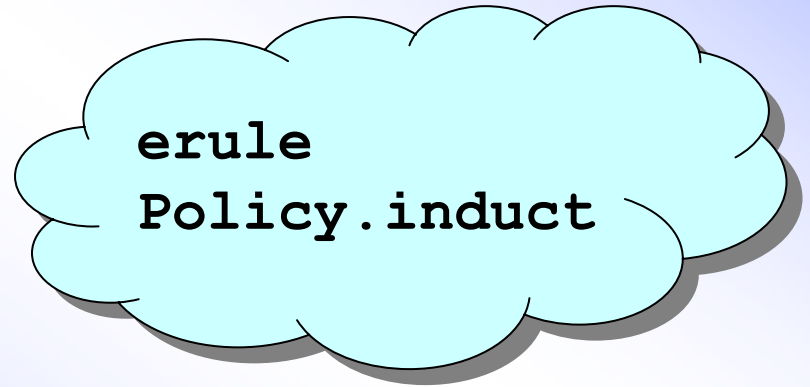
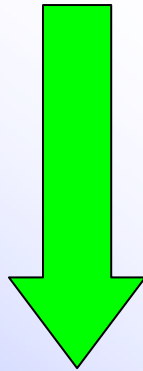
theorem No_Dilemma :

“ [| nm ∈ set tr; tr ∈ Policy |] ⇒
Dilemma nm ∉ set tr ”

Proof - step 1

`tr ∈ Policy ⇒`

`nm ∈ set tr → Contradiction nm → set tr`



12 subgoal!!!


Proof – step 2

[|tr1 ∈ Policy; play a User; public f;

nm ∈ set tr1 → Contradiction nm ∉ set tr1|]

⇒ nm ∈ set (Perm (Read a f) # tr1) →

Contradiction nm ∉ set (Perm(Read a f) # tr1)



simp del: "Perm_Obl"

Proof – step 3

[| tr1 ∈ Policy; play a User; public f;

nm ∈ set tr1 → Contradiction nm ∉ set tr1 |]

⇒ (nm = Perm (Read a f)) →

Contradiction (Perm (Read a f)) ≠ Perm (Read a f) ∧

Contradiction (Perm (Read a f)) ∉ set tr1)

∧ (nm ∈ set tr1 → Contradiction nm ≠ Perm(Read a f))



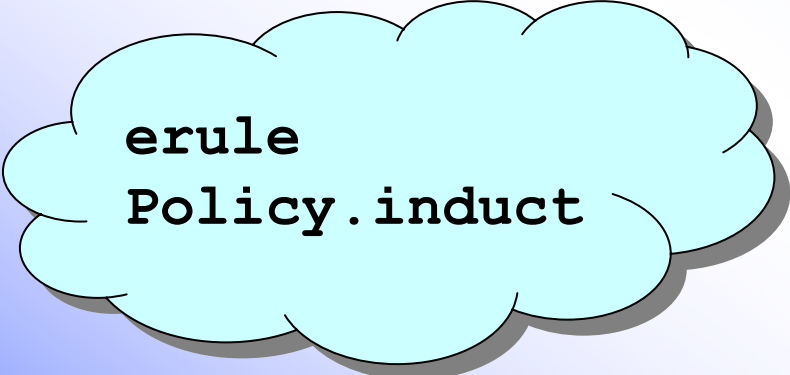
subgoal_tac
"Contradiction(Perm(Read a f))
∉ set tr1"

Proof – step 4

[| tr1 ∈ Policy; play a User; public f;

nm ∈ set tr1 → Contradiction nm ∉ set tr1 |]

⇒ Contradiction (Perm (Read a f)) ∉ set tr1



erule

Policy.induct

Proof – step 5

[| tr1 ∈ Policy; play a User; public f;

nm ∈ set tr1 → Contradiction nm ∉ set tr1;

Contradiction (Perm (Read a f)) ∉ set tr1|]

⇒ (nm = Perm(Read a f) →

Contradiction(Perm(Read a f)) ≠ Perm(Read a f) ∧

Contradiction(Perm(Read a f)) ∉ set tr1)

∧ (nm ∈ set tr1 → Contradiction nm ≠ Perm(Read a f))



rule conjI

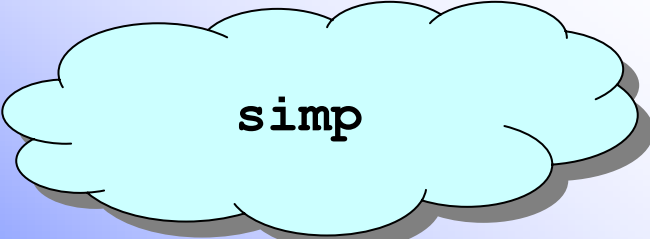
Proof – step 6

[| tr1 ∈ Policy; play a User; public f;
nm ∈ set tr1 → Contradiction nm ∉ set tr1;
Contradiction(Perm(Read a f)) ∉ set tr1|]

⇒ nm = Perm(Read a f) →

Contradiction(Perm(Read a f)) ≠ Perm(Read a f)

∧ Contradiction(Perm(Read a f)) ∉ set tr1



simp

Proof – step 7

[| tr1 ∈ Policy; play a User; public f;

nm ∈ set tr1 → Contradiction nm ∉ set tr1;

Contradiction (Perm (Read a f)) ∉ set tr1|]

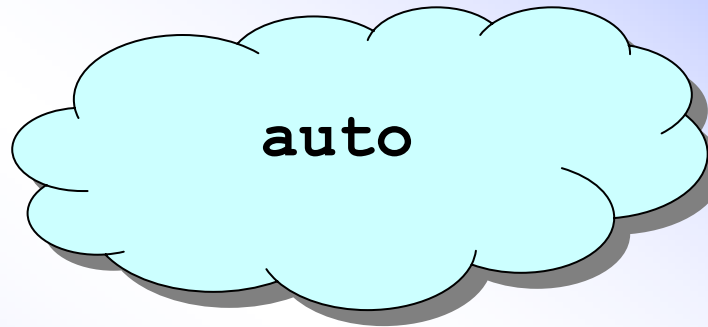
⇒ nm ∈ set tr1 → Contradiction nm ≠ Perm(Read a f)



erule

Policy.induct

Proof – step 8



```
 $\Lambda$  a f tr8 tr1.[| play a Bad; tr1  $\in$  Policy;  
public f;  $\neg$ n (Obl ( $\neg$ o (Read a f)))  $\notin$  set tr8;  
 $\neg$ n (Obl ( $\neg$ o (Read a f)))  $\notin$  set tr1|]  
 $\Rightarrow$  False
```

Policy Inconsistencies

6 Contradictions:

- N7 – N3:** “A system security officer is both permitted and forbidden to downgrade a public file”
- N8 – N1:** “A bad user is both forbidden and permitted to read a public file”
- N8 – N5:** “A bad user is both forbidden and permitted to read a not public file”
- N9 – N2:** “A bad user is both forbidden and permitted to write on a public file he owns”
- N9 – N6:** “A bad user is both forbidden and permitted to write on a not public file he owns”
- N10 – N7:** “A bad user is both forbidden and permitted to downgrade a file”

1 Dilemma:

- N11 – N4 :** “A bad user is both forbidden and obliged to change his password”

Conclusions

- Developed the first inductive approach to prove security policy correctness
- Mechanized the approach with the proof assistant *Isabelle*
- Verified presence of many inconsistencies in the example policy: proof script of 500 lines

Next steps...

- To simplify proof demonstration strategy
- Search of alternative formalization, if possible without trace
- Application to widest study case
- Extension to union of more policy



Thanks!!!